

Realizability and Dialectica

(with partial functions)

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Proofs and Computation

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A TYPE-FREE GÖDEL INTERPRETATION

MICHAEL BEESON

The essential feature by which the present interpretation can be distinguished from Gödel's is this: Gödel's interpretation makes essential use of functionals of higher type. A functional of type (σ, \mathfrak{s}) must be defined on all type σ functionals and take type \mathfrak{s} functionals for values. Our interpretation, however, uses partially defined functions (better: operations).

We have thus solved the proportion:

X : Gödel's interpretation as

Kleene's realizability : Kreisel's "modified realizability".

Plan

- **Realizability** (*formulas as sets*)
- **Dialectica** (*formulas as relations*)
- **Point 1:** Relational view is more general
- **Point 2:** Thoughts on Beeson's question

Realizability

Kleene Realizability

$$\underbrace{A}_{\text{sentence}} \mapsto \underbrace{\{n : n r A\}}_{\text{set of realisers of } A}$$

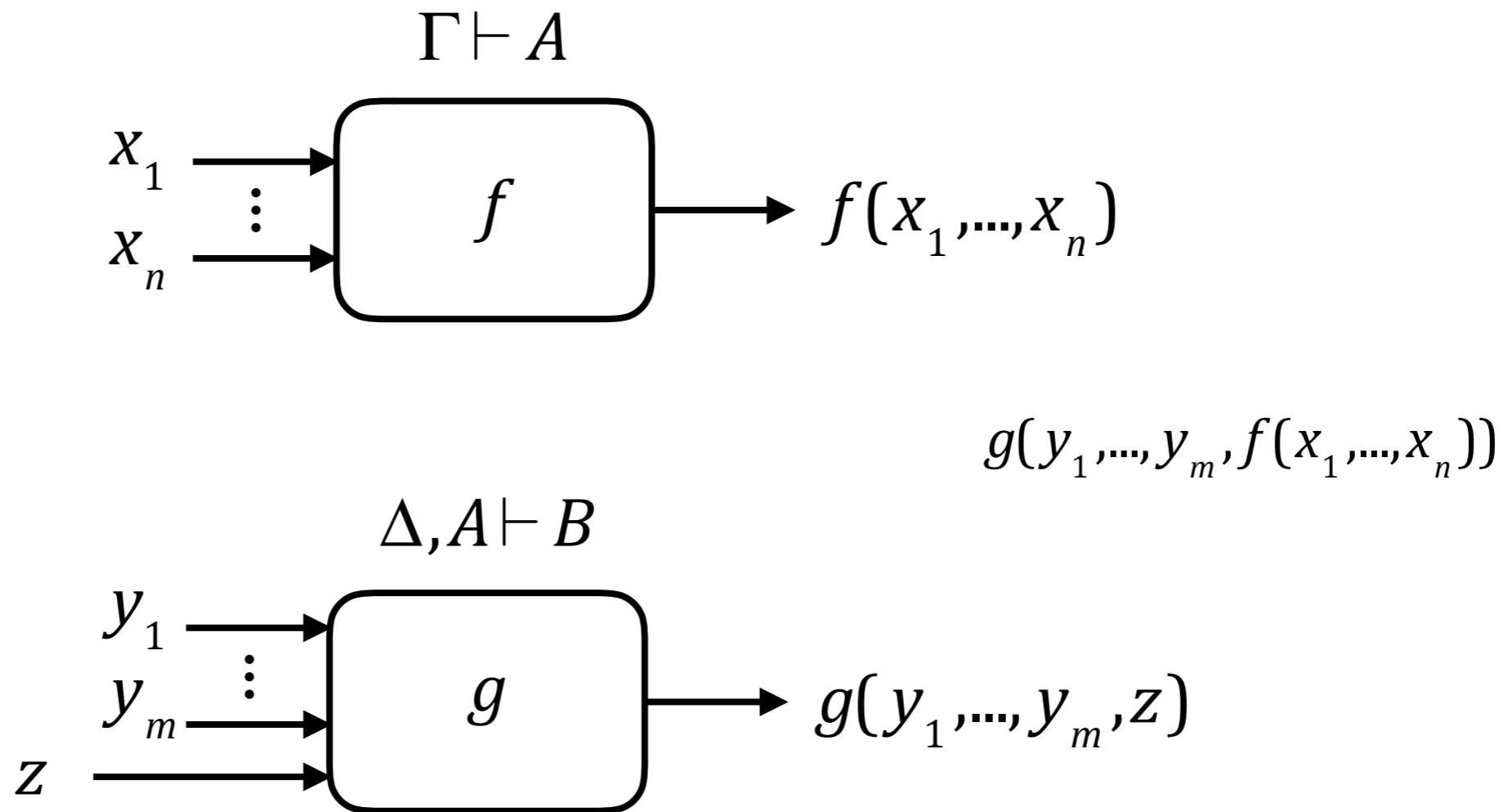
ON THE INTERPRETATION OF INTUITIONISTIC NUMBER THEORY

S. C. KLEENE

$$\begin{aligned}e \text{ r } F &\equiv e = 0 \wedge F \\e \text{ r } A \wedge B &\equiv e_0 \text{ r } A \wedge e_1 \text{ r } B \\e \text{ r } A \vee B &\equiv (e_0 = 0 \wedge e_1 \text{ r } A) \wedge (e_0 = 1 \wedge e_1 \text{ r } B) \\e \text{ r } A \rightarrow B &\equiv \forall a(a \text{ r } A \rightarrow \underline{\{e\}(a)} \downarrow \wedge \{e\}(a) \text{ r } B) \\e \text{ r } \exists x A(x) &\equiv e_1 \text{ r } A(e_0) \\e \text{ r } \forall x A(x) &\equiv \forall x(\{e\}(x) \downarrow \wedge \{e\}(x) \text{ r } A(x))\end{aligned}$$

Theorem (Kleene'45). If $HA \vdash A$ then there exists a numeral n such that $HA \vdash n \text{ r } A$

Kleene Realizability



Dialectica

Dialectica Interpretation

$$\underbrace{A}_{\text{sentence}} \quad \mapsto \quad \{ (x, y) : \underbrace{|A|_y^x}_{\text{relation between arguments and counter-arguments}} \}$$

ÜBER EINE BISHER NOCH NICHT BENÜTZTE ERWEITERUNG DES FINITEN STANDPUNKTES

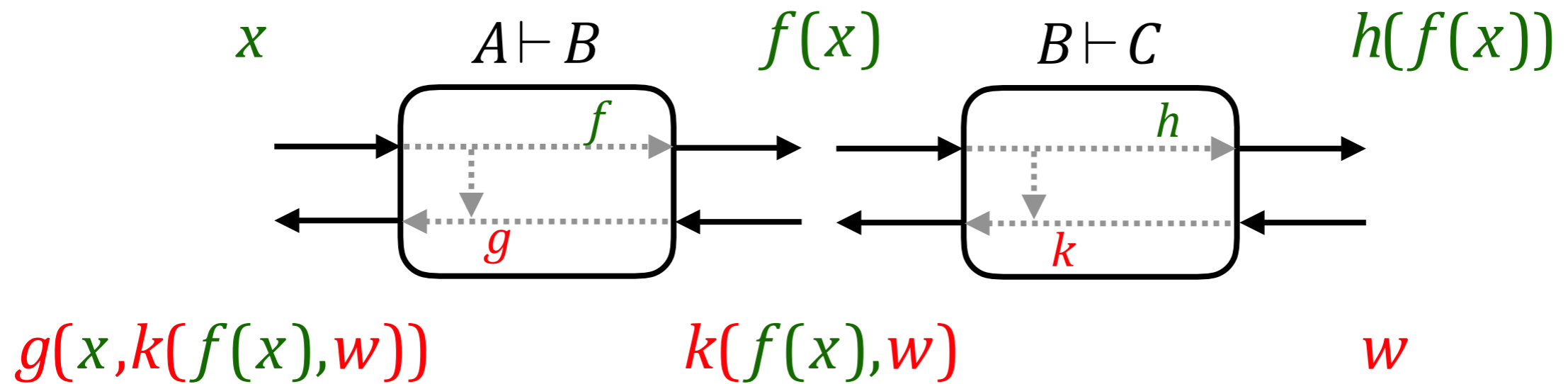
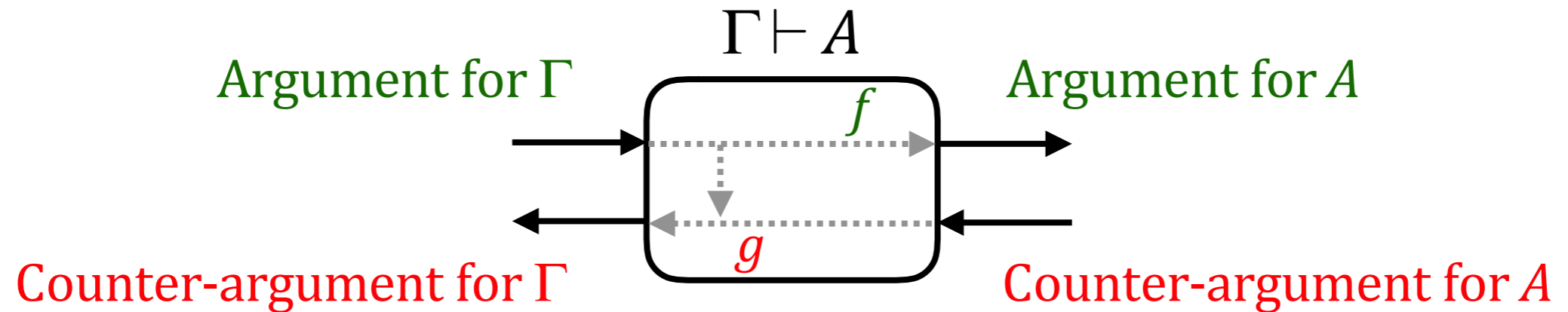
von Kurt GÖDEL, Princeton

Dialectica, vol. 12, 1958

$$\begin{aligned} |F \wedge G|_{z,w}^{y,v} &\equiv |F|_z^y \wedge |G|_w^v \\ |F \vee G|_{z,w}^{y,v,t} &\equiv (t=0 \wedge |F|_z^y) \vee (t=1 \wedge |G|_w^v) \\ |F \rightarrow G|_{y,w}^{v,z} &\equiv |F|_{z(y,w)}^y \rightarrow |G|_w^{v(y)} \\ |\forall x F(x)|_{s,z}^y &\equiv |F(s)|_z^{y(s)} \\ |\exists x F(x)|_z^{s,y} &\equiv |F(s)|_z^y \end{aligned}$$

Theorem (Gödel'58). If $HA \vdash A$ then there exists a term t of system T such that $T \vdash \forall y |A|_y^t$

Dialectica Interpretation



Point 1:
Kleene realizability
in 'relational' style



Gödel's Dialectica interpretation

$$\begin{aligned}
 |F \wedge G|_{z,w}^{y,v} &\equiv |F|_z^y \wedge |G|_w^v \\
 |F \vee G|_{z,w}^{y,v,t} &\equiv (t=0 \wedge |F|_z^y) \vee (t=1 \wedge |G|_w^v) \\
 |F \rightarrow G|_{y,w}^{v,z} &\equiv |F|_{Z(y,w)}^y \rightarrow |G|_w^{V(y)} \\
 |\forall s F|_{s,z}^y &\equiv |F|_z^{Y(s)} \\
 |\exists s F|_z^{s,y} &\equiv |F|_z^y
 \end{aligned}$$

Kreisel's modified realizability

$$\begin{aligned}
 |F \wedge G|_{z,w}^{y,v} &\equiv |F|_z^y \wedge |G|_w^v \\
 |F \vee G|_{z,w}^{y,v,t} &\equiv (t=0 \wedge |F|_z^y) \vee (t=1 \wedge |G|_w^v) \\
 |F \rightarrow G|_{y,w}^v &\equiv \forall z |F|_z^y \rightarrow |G|_w^{V(y)} \\
 |\forall s F|_{s,z}^y &\equiv |F|_z^{Y(s)} \\
 |\exists s F|_z^{s,y} &\equiv |F|_z^y
 \end{aligned}$$

$$A^D(x, y) \text{ iff } |A|_y^x$$

$$x \text{ mr } A \text{ iff } \forall y |A|_y^x$$

Kleene
numerical
realizability

$$\begin{aligned}
 e \text{ r } A \wedge B &\equiv e_0 \text{ r } A \wedge e_1 \text{ r } B \\
 e \text{ r } A \vee B &\equiv (e_0 = 0 \wedge e_1 \text{ r } A) \wedge (e_0 = 1 \wedge e_1 \text{ r } B) \\
 e \text{ r } A \rightarrow B &\equiv \forall a (a \text{ r } A \rightarrow \{e\}(a) \downarrow \wedge \{e\}(a) \text{ r } B) \\
 e \text{ r } \exists x A(x) &\equiv e_1 \text{ r } A(e_0) \\
 e \text{ r } \forall x A(x) &\equiv \forall x (\{e\}(x) \downarrow \wedge \{e\}(x) \text{ r } A(x))
 \end{aligned}$$

Kleene (*dialog*)
numerical
realizability

$$\begin{aligned}
 |F \wedge G|_a^e &\equiv |F|_{a_0}^{e_0} \wedge |G|_{a_1}^{e_1} \\
 |F \vee G|_a^e &\equiv (e_0 = 0 \wedge |F|_a^{e_1}) \vee (e_0 = 1 \wedge |G|_a^{e_1}) \\
 |F \rightarrow G|_a^e &\equiv \forall b |F|_b^{a_0} \rightarrow \{e\}(a_0) \downarrow \wedge |G|_{a_1}^{\{e\}(a_0)} \\
 |\exists s F(s)|_a^e &\equiv |F(e_0)|_a^{e_1} \\
 |\forall s F(s)|_a^e &\equiv \{e\}(a_0) \downarrow \wedge |F(a_0)|_{a_1}^{\{e\}(a_0)}
 \end{aligned}$$

$$e \text{ r } A \text{ iff } \forall a |A|_a^e$$

Point 2:
Dialectica with
partial realizers

Kleene realizability

$$|F \rightarrow G|_a^e \equiv \forall b |F|_b^{a_0} \rightarrow \{e\}(a_0) \downarrow \wedge |G|_{a_1}^{\{e\}(a_0)}$$

Attempt 1 (problem interpreting contraction axiom)

$$|F \rightarrow G|_a^e \equiv (\{e_1\}(a) \downarrow \rightarrow |F|_{\{e_1\}(a)}^{a_0}) \rightarrow (\{e_0\}(a_0) \downarrow \wedge |G|_{a_1}^{\{e_0\}(a_0)})$$

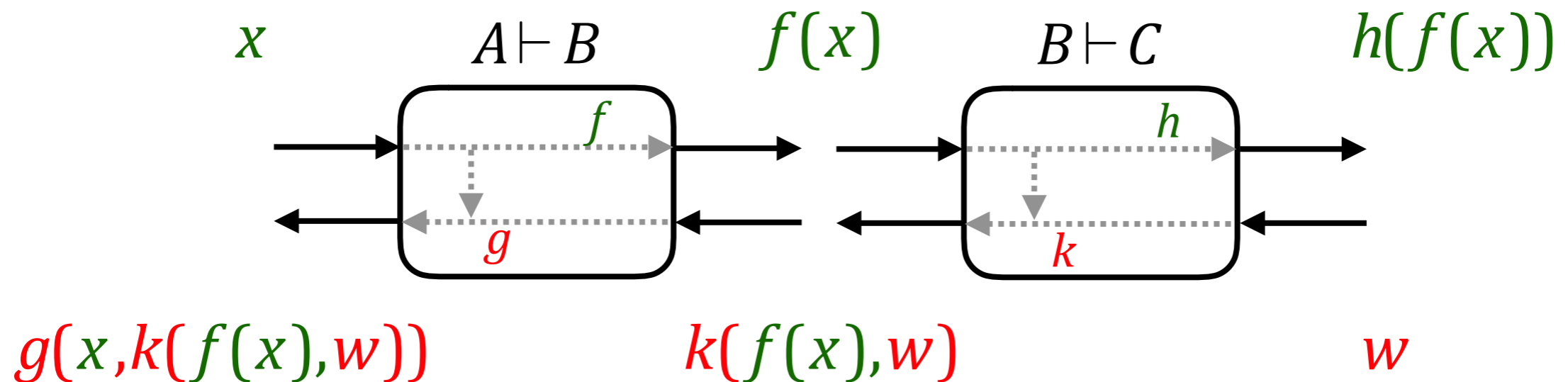
Attempt 2 (still has issue with contraction axiom)

$$|F \rightarrow G|_a^e \equiv (\{e_1\}(a) \downarrow \rightarrow \forall n \in \{e_1\}(a) |F|_n^{a_0}) \rightarrow \{e_0\}(a_0) \downarrow \wedge |G|_{a_1}^{\{e_0\}(a_0)}$$

Attempt 3 (problem interpreting composition)

$$|F \rightarrow G|_a^e \equiv \{e_1\}(a) \downarrow \wedge (\forall n \in \{e_1\}(a) |F|_n^{a_0} \rightarrow \{e_0\}(a_0) \downarrow \wedge |G|_{a_1}^{\{e_0\}(a_0)})$$

Issue with Composition



If **arguments** are allowed to be partial, we can't guarantee **counter-arguments** are always total

Possible Solution

We need a finite set $W_{\{g\}(x,w)}$ so that we can define:

$$|F \rightarrow G|_{x,w}^{f,g} \equiv \forall y \in W_{\{g\}(x,w)} |F|_y^x \rightarrow \{f\}(x) \downarrow \wedge |G|_w^{\{f\}(x)}$$

But we need to be able to:

- form finite union of W -sets (*contraction*)
- form union of W -sets indexed by W -sets (*composition*)
- W -sets need to be non-empty (*composition*)

Possible Solution

Define W -sets as:

$$W_{e,x} \equiv \{0\} \cup \{ \{e\}(i,x) \mid i \in \mathbb{N} \wedge \{e\}(i,x) \downarrow \}$$

So:

$$\forall y \in W_{e,x} A(y) \equiv A(0) \wedge \forall i,u (T(e,(i,x),u) \rightarrow A(U(u)))$$

Summary

- Kleene realizability in “relational style”
- When trying to witness counter-argument function
 - Can’t witness it precisely because matrix not decidable
 - Can’t assume counter-argument functions are total
 - Defined finite sets of possibly undefined elements
 - These sets must be non-empty